JFP **13** (5): 957–959, September 2003. © 2003 Cambridge University Press
DOI: 10.1017/S0956796802004446 Printed in the United Kingdom

ERRATA

Stack-based typed assembly language

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The following three figures (figures 10, 11 and 12) were not shown in the original published version of the article. These figures constitute the entire static semantics of the STAL type system.

$$\boxed{\Delta \vdash \tau \quad \Delta \vdash \sigma \quad \vdash \Psi \quad \Delta \vdash \Gamma}$$

$$(\text{type})\ \frac{}{\Delta \vdash \tau}\ (\text{fv}(\tau) \subseteq \Delta) \qquad (\text{stype})\ \frac{}{\Delta \vdash \sigma}\ (\text{fv}(\sigma) \subseteq \Delta)$$

$$(\text{htype})\ \frac{\cdot \vdash \tau_1 \quad \cdots \quad \cdot \vdash \tau_n}{\vdash \{\ell_1{:}\tau_1, \ldots, \ell_n{:}\tau_n\}} \qquad (\text{rftype})\ \frac{\Delta \vdash \sigma \quad \Delta \vdash \tau_1 \quad \cdots \quad \Delta \vdash \tau_n}{\Delta \vdash \{\mathtt{sp}{:}\sigma, r_1{:}\tau_1, \ldots, r_n{:}\tau_n\}}$$

$$\boxed{\Delta \vdash \sigma_1 = \sigma_2}$$

$$(\text{seq-refl})\ \frac{\Delta \vdash \sigma}{\Delta \vdash \sigma = \sigma} \qquad (\text{seq-sym})\ \frac{\Delta \vdash \sigma_2 = \sigma_1}{\Delta \vdash \sigma_1 = \sigma_2}$$

$$(\text{seq-trans})\ \frac{\Delta \vdash \sigma_1 = \sigma_2 \quad \Delta \vdash \sigma_2 = \sigma_3}{\Delta \vdash \sigma_1 = \sigma_3}$$

$$(\text{seq-cons})\ \frac{\Delta \vdash \tau \quad \Delta \vdash \sigma_1 = \sigma_2}{\Delta \vdash \tau{::}\sigma_1 = \tau{::}\sigma_2} \qquad (\text{seq-append})\ \frac{\Delta \vdash \sigma_1 = \sigma_1' \quad \Delta \vdash \sigma_2 = \sigma_2'}{\Delta \vdash \sigma_1 \mathbin{@} \sigma_2 = \sigma_1' \mathbin{@} \sigma_2'}$$

$$(\text{stk}\beta 1)\ \frac{\Delta \vdash \sigma}{\Delta \vdash \mathit{nil} \mathbin{@} \sigma = \sigma} \qquad (\text{stk}\beta 2)\ \frac{\Delta \vdash \sigma}{\Delta \vdash \sigma \mathbin{@} \mathit{nil} = \sigma}$$

$$(\text{stk}\beta 3)\ \frac{\Delta \vdash \tau \quad \Delta \vdash \sigma_1 \quad \Delta \vdash \sigma_2}{\Delta \vdash (\tau{::}\sigma_1) \mathbin{@} \sigma_2 = \tau{::}(\sigma_1 \mathbin{@} \sigma_2)}$$

$$(\text{stk}\beta 4)\ \frac{\Delta \vdash \sigma_1 \quad \Delta \vdash \sigma_2 \quad \Delta \vdash \sigma_3}{\Delta \vdash (\sigma_1 \mathbin{@} \sigma_2) \mathbin{@} \sigma_3 = \sigma_1 \mathbin{@} (\sigma_2 \mathbin{@} \sigma_3)}$$

$$\boxed{\Delta \vdash \Gamma_1 \leq \Gamma_2}$$

$$(\text{rf-leq})\ \frac{\Delta \vdash \tau_i \quad (\text{for } 1 \leq i \leq m) \quad \Delta \vdash \sigma = \sigma'}{\Delta \vdash \{\mathtt{sp}{:}\sigma, r_1{:}\tau_1, \ldots, r_m{:}\tau_m\} \leq \{\mathtt{sp}{:}\sigma', r_1{:}\tau_1, \ldots, r_n{:}\tau_n\}}\ (m \geq n)$$

Fig. 10. Static Semantics of STAL, Judgments for Types

$$\boxed{\vdash M \quad \vdash H : \Psi \quad \Psi \vdash S : \sigma \quad \Psi \vdash R : \Gamma}$$

$$(\text{mach})\ \frac{\vdash H : \Psi \quad \Psi \vdash R : \Gamma \quad \Psi; \cdot; \Gamma \vdash I}{\vdash (H, R, I)}$$

$$(\text{heap})\ \frac{\vdash \Psi \quad \Psi \vdash h_i : \Psi(\ell_i)\ \text{hval} \quad (\text{for } 1 \le i \le n)}{\vdash \{\ell_1 \mapsto h_1, \ldots, \ell_n \mapsto h_n\} : \Psi}$$

$$(\text{nil})\ \frac{}{\Psi \vdash nil : nil} \qquad (\text{cons})\ \frac{\Psi; \cdot; \cdot \vdash w : \tau \quad \Psi \vdash S : \sigma}{\Psi \vdash w{::}S : \tau{::}\sigma} \qquad (\text{stkeq})\ \frac{\Psi \vdash S : \sigma_1 \quad \cdot \vdash \sigma_1 = \sigma_2}{\Psi \vdash S : \sigma_2}$$

$$(\text{regfile})\ \frac{\Psi \vdash S : \sigma \quad \Psi; \cdot; \cdot \vdash w_i : \tau_i \quad (\text{for } 1 \le i \le n)}{\Psi \vdash \{\text{sp} \mapsto S, r_1 \mapsto w_1, \ldots, r_m \mapsto w_m\} : \{\text{sp}{:}\sigma, r_1{:}\tau_1, \ldots, r_n{:}\tau_n\}}\ (m \ge n)$$

$$\boxed{\Psi \vdash h : \tau\ \text{hval} \quad \Psi; \Delta; \Gamma \vdash v : \tau}$$

$$(\text{tuple})\ \frac{\Psi; \cdot; \cdot \vdash w_i : \tau_i}{\Psi \vdash \langle w_1, \ldots, w_n \rangle : \langle \tau_1, \ldots, \tau_n \rangle\ \text{hval}} \qquad (\text{code})\ \frac{\Delta \vdash \Gamma \quad \Psi; \Delta; \Gamma \vdash I}{\Psi \vdash \texttt{code}[\Delta]\Gamma.I : \forall[\Delta].\Gamma\ \text{hval}}$$

$$(\text{label})\ \frac{}{\Psi; \Delta; \Gamma \vdash \ell : \Psi(\ell)} \qquad (\text{int})\ \frac{}{\Psi; \Delta; \Gamma \vdash i : int}$$

$$(\text{ns})\ \frac{}{\Psi; \Delta; \Gamma \vdash ns : \top} \qquad (\text{ptr})\ \frac{\Delta \vdash \sigma}{\Psi; \Delta; \Gamma \vdash ptr(i) : ptr(\sigma)}\ (|\sigma| = i)$$

$$(\text{reg})\ \frac{}{\Psi; \Delta; \Gamma \vdash r : \Gamma(r)}$$

$$(\text{tapp})\ \frac{\Delta \vdash \tau \quad \Psi; \Delta; \Gamma \vdash v : \forall[\alpha, \Delta'].\Gamma'}{\Psi; \Delta; \Gamma \vdash v[\tau] : \forall[\Delta'].\Gamma'[\tau/\alpha]} \qquad (\text{stapp})\ \frac{\Delta \vdash \sigma \quad \Psi; \Delta; \Gamma \vdash v : \forall[\rho, \Delta'].\Gamma'}{\Psi; \Delta; \Gamma \vdash v[\sigma] : \forall[\Delta'].\Gamma'[\sigma/\rho]}$$

$$(\text{pack})\ \frac{\Delta \vdash \tau \quad \Psi; \Delta; \Gamma \vdash v : \tau'[\tau/\alpha]}{\Psi; \Delta; \Gamma \vdash pack\ [\tau, v]\ as\ \exists\alpha.\tau' : \exists\alpha.\tau'}$$

$$\boxed{\Psi; \Delta; \Gamma \vdash I}$$

$$(\text{seq})\ \frac{\Psi; \Delta; \Gamma \vdash \iota \Rightarrow \Delta'; \Gamma' \quad \Psi; \Delta'; \Gamma' \vdash I}{\Psi; \Delta; \Gamma \vdash \iota; I} \qquad (\text{jmp})\ \frac{\Delta \vdash \Gamma_1 \le \Gamma_2 \quad \Psi; \Delta; \Gamma_1 \vdash v : \forall[\,].\Gamma_2}{\Psi; \Delta; \Gamma_1 \vdash \texttt{jmp}\ v}$$

$$(\text{halt})\ \frac{\Psi; \Delta; \Gamma \vdash \texttt{r1} : \tau}{\Psi; \Delta; \Gamma \vdash \texttt{halt}[\tau]}$$

Fig. 11. STAL Static Semantics, Term Constructs except Instructions

https://doi.org/10.1017/S0956796802004446 Published online by Cambridge University Press

$$(\text{aop})\ \frac{\Psi;\Delta;\Gamma\vdash r_s : \mathit{int} \quad \Psi;\Delta;\Gamma\vdash v : \mathit{int}}{\Psi;\Delta;\Gamma\vdash \mathit{aop}\ \ r_d, r_s, v \Rightarrow \Delta;\Gamma\{r_d{:}\mathit{int}\}}$$

$$(\text{bop})\ \frac{\Psi;\Delta;\Gamma_1\vdash r : \mathit{int} \quad \Psi;\Delta;\Gamma_1\vdash v : \forall[\,].\Gamma_2 \quad \Delta\vdash\Gamma_1\leq\Gamma_2}{\Psi;\Delta;\Gamma_1\vdash \mathit{bop}\ \ r, v \Rightarrow \Delta;\Gamma_1}$$

$$(\text{ld})\ \frac{\Psi;\Delta;\Gamma\vdash r_s : \langle\tau_0,\ldots,\tau_{n-1}\rangle}{\Psi;\Delta;\Gamma\vdash \mathtt{ld}\ \ r_d, r_s(i) \Rightarrow \Delta;\Gamma\{r_d{:}\tau_i\}}\ (0\leq i<n)$$

$$(\text{malloc})\ \frac{\Psi;\Delta;\Gamma\vdash v_i : \tau_i}{\Psi;\Delta;\Gamma\vdash \mathtt{malloc}\ \ r_d, \langle v_1,\ldots,v_n\rangle \Rightarrow \Delta;\Gamma\{r_d{:}\langle\tau_1,\ldots,\tau_n\rangle\}}\ (1\leq i\leq n)$$

$$(\text{mov})\ \frac{\Psi;\Delta;\Gamma\vdash v : \tau}{\Psi;\Delta;\Gamma\vdash \mathtt{mov}\ \ r_d, v \Rightarrow \Delta;\Gamma\{r_d{:}\tau\}}$$

$$(\text{unpack})\ \frac{\Psi;\Delta;\Gamma\vdash v : \exists\alpha.\tau}{\Psi;\Delta;\Gamma\vdash \mathtt{unpack}\ \ [\alpha, r_d], v \Rightarrow \alpha,\Delta;\Gamma\{r_d{:}\tau\}}\ (\alpha\notin\Delta)$$

$$(\text{get-sp})\ \frac{}{\Psi;\Delta;\Gamma\vdash \mathtt{mov}\ \ r_d, \mathtt{sp} \Rightarrow \Delta;\Gamma\{r_d{:}\mathit{ptr}(\sigma)\}}\ (\Gamma(\mathtt{sp})=\sigma)$$

$$(\text{set-sp})\ \frac{\Psi;\Delta;\Gamma\vdash r_s : \mathit{ptr}(\sigma_2) \quad \Delta\vdash\sigma_1=\sigma_3\mathbin{@}\sigma_2}{\Psi;\Delta;\Gamma\vdash \mathtt{mov}\ \ \mathtt{sp}, r_s \Rightarrow \Delta;\Gamma\{\mathtt{sp}{:}\sigma_2\}}\ (\Gamma(\mathtt{sp})=\sigma_1)$$

$$(\text{salloc})\ \frac{}{\Psi;\Delta;\Gamma\vdash \mathtt{salloc}\ \ n \Rightarrow \Delta;\Gamma\{\mathtt{sp}{:}\underbrace{\top::\cdots::\top}_{n}::\sigma\}}\ (\Gamma(\mathtt{sp})=\sigma)$$

$$(\text{sfree})\ \frac{\Delta\vdash\sigma_1=\tau_0::\cdots::\tau_{n-1}::\sigma_2}{\Psi;\Delta;\Gamma\vdash \mathtt{sfree}\ \ n \Rightarrow \Delta;\Gamma\{\mathtt{sp}{:}\sigma_2\}}\ (\Gamma(\mathtt{sp})=\sigma_1)$$

$$(\text{sld1})\ \frac{\Delta\vdash\sigma_1=\tau_0::\cdots::\tau_i::\sigma_2}{\Psi;\Delta;\Gamma\vdash \mathtt{sld}\ \ r_d, \mathtt{sp}(i) \Rightarrow \Delta;\Gamma\{r_d{:}\tau_i\}}\ (\Gamma(\mathtt{sp})=\sigma_1\wedge 0\leq i)$$

$$(\text{sld2})\ \frac{\begin{array}{c}\Psi;\Delta;\Gamma\vdash r_s : \mathit{ptr}(\sigma_3) \qquad \Delta\vdash\sigma_1=\sigma_2\mathbin{@}\sigma_3\\ \Delta\vdash\sigma_3=\tau_0::\cdots::\tau_i::\sigma_4\end{array}}{\Psi;\Delta;\Gamma\vdash \mathtt{sld}\ \ r_d, r_s(i) \Rightarrow \Delta;\Gamma\{r_d{:}\tau_i\}}\ (\Gamma(\mathtt{sp})=\sigma_1\wedge 0\leq i)$$

$$(\text{sst1})\ \frac{\Delta\vdash\sigma_1=\tau_0::\cdots::\tau_i::\sigma_2 \quad \Psi;\Delta;\Gamma\vdash r_s : \tau}{\Psi;\Delta;\Gamma\vdash \mathtt{sst}\ \ \mathtt{sp}(i), r_s \Rightarrow \Delta;\Gamma\{\mathtt{sp}{:}\tau_0::\cdots::\tau_{i-1}::\tau::\sigma_2\}}\ (\Gamma(\mathtt{sp})=\sigma_1\wedge 0\leq i)$$

$$(\text{sst2})\ \frac{\begin{array}{c}\Psi;\Delta;\Gamma\vdash r_d : \mathit{ptr}(\sigma_3) \qquad \Psi;\Delta;\Gamma\vdash r_s : \tau\\ \Delta\vdash\sigma_1=\sigma_2\mathbin{@}\sigma_3 \qquad \Delta\vdash\sigma_3=\tau_0::\cdots::\tau_i::\sigma_4\\ \Delta\vdash\sigma_5=\tau_0::\cdots::\tau_{i-1}::\tau::\sigma_4\end{array}}{\Psi;\Delta;\Gamma\vdash \mathtt{sst}\ \ r_d(i), r_s \Rightarrow \Delta;\Gamma\{\mathtt{sp}{:}\sigma_2\mathbin{@}\sigma_5, r_d{:}\mathit{ptr}(\sigma_5)\}}\ (\Gamma(\mathtt{sp})=\sigma_1\wedge 0\leq i)$$

$$(\text{st})\ \frac{\Psi;\Delta;\Gamma\vdash r_d : \langle\tau_0,\ldots,\tau_{n-1}\rangle \quad \Psi;\Delta;\Gamma\vdash r_s : \tau_i}{\Psi;\Delta;\Gamma\vdash \mathtt{st}\ \ r_d(i), r_s \Rightarrow \Delta;\Gamma}\ (0\leq i<n)$$

Fig. 12. STAL Static Semantics, Instructions

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