

TOURNAMENTS WHOSE SUBTOURNAMENTS ARE IRREDUCIBLE OR TRANSITIVE

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ABSTRACT. Beineke and Harary gave an example of a family of tournaments T_n such that every subtournament of T_n is irreducible or transitive. We characterize all tournaments with this property.

1. Introduction. A tournament T_n consists of a finite set of nodes $1, 2, \dots, n$ such that each pair of distinct nodes i and j is joined by exactly one of the arcs \overrightarrow{ij} or \overrightarrow{ji} . If the arc \overrightarrow{ij} is in T_n we say that i beats j or j loses to i and write $i \rightarrow j$. If each node of a subtournament A beats each node of a subtournament B , we write $A \rightarrow B$ and let $A + B$ denote the subtournament determined by the nodes of A and B .

A tournament T_n is *reducible* if it can be expressed as $T_n = A + B$ for some non-empty tournaments A and B ; otherwise it is *irreducible*. A tournament is *transitive* if there exists a linear ordering of its nodes such that $i \rightarrow j$ if and only if i precedes j in the ordering. (Notice that the trivial tournament T_1 is the only tournament that is both transitive and irreducible.) A tournament T_n is *highly-regular* if n is odd and there exists a cyclic ordering of the nodes such that $i \rightarrow j$ if and only if j is one of the first $\frac{1}{2}(n-1)$ successors of i in the ordering; we remark that the ordering with this property is unique. (These tournaments were introduced by Kendall and Babington Smith [3]. For additional material on tournaments in general, see [2] and [4].)

We say a tournament T_n has *property \mathcal{L}* if every subtournament of T_n is irreducible or transitive. Beineke and Harary [1] showed that highly-regular tournaments have property \mathcal{L} . Our main object here is to establish a structural characterization of all tournaments with property \mathcal{L} . Before stating this characterization we need to introduce some additional terminology.

2. Statement of characterization. We say the nodes in a subtournament A of T_n are *equivalent* if for any node q not in A either $q \rightarrow A$ or $A \rightarrow q$. (Equivalent nodes are sometimes said to form a *convex* subset; see, e.g., [6]). Suppose the nodes of T_n are partitioned into disjoint subtournaments E_1, \dots, E_m of equivalent nodes, where the subscripts $1, \dots, m$ serve merely to distinguish between different subtournaments. Then $E_i \rightarrow E_j$ or $E_j \rightarrow E_i$ for

$1 \leq i, j \leq m$. If R_m denotes the tournament on m nodes in which $i \rightarrow j$ if and only if $E_i \rightarrow E_j$, then we write $T_n = R_m(E_1, \dots, E_m)$. (If the subtournaments E_1, \dots, E_m are isomorphic, then T_n is the *composition* of R_m with E_1 ; see [4; p.78].) A tournament T_n is *simple* (see, e.g., [2] and [6]) if it has no non-trivial proper subtournaments of equivalent nodes, that is, if the equation $T_n = R_m(E_1, \dots, E_m)$ implies that $m = 1$ and $E_1 = T_n$ or that $m = n$, $T_n = R_m$, and $E_i = T_1$ for each i . We can now state our main result.

THEOREM 1. *A tournament T_n has property \mathcal{L} if and only if*

$$T_n = R_m(E_1, \dots, E_m)$$

where R_m is a highly-regular tournament and the subtournaments E_1, \dots, E_m all are transitive.

3. Two preliminary results. For any node i in a tournament T_n , let $\Gamma(i)$ and $\Gamma^{-1}(i)$ denote the subtournaments determined by the nodes of T_n that lose to i and the nodes that beat i , respectively. It may be that $\Gamma(i)$ or $\Gamma^{-1}(i)$ is the empty tournament. We shall use the following lemmas in the proof of Theorem 1.

LEMMA 1. *A tournament T_n has property \mathcal{L} if and only if $\Gamma(i)$ and $\Gamma^{-1}(i)$ are empty or transitive for all nodes i of T_n .*

Proof. If $\Gamma(i)$ or $\Gamma^{-1}(i)$ is non-empty and non-transitive for some node i of T_n , then the corresponding subtournament $i + \Gamma(i)$ or $\Gamma^{-1}(i) + i$ is neither irreducible nor transitive and, consequently, T_n does not have property \mathcal{L} . Conversely, if T_n does not have property \mathcal{L} , then it contains a reducible non-transitive subtournament $S = A + B$ where A and B are both non-empty and at least one of them is non-transitive. Let i and j denote any nodes in A and B , respectively. If B is non-transitive then $\Gamma(i)$ is neither empty nor transitive, and if A is non-transitive then $\Gamma^{-1}(j)$ is neither empty nor transitive. This completes the proof of Lemma 1.

LEMMA 2. *A non-simple tournament $T_n = R_m(E_1, \dots, E_m)$, where $1 < m < n$, has property \mathcal{L} if and only if the tournament R_m has property \mathcal{L} and the subtournaments E_1, \dots, E_m all are transitive.*

Proof. The sufficiency of the conditions follows readily from Lemma 1 and the necessity of the condition that R_m must have property \mathcal{L} is obvious. Suppose some subtournament E_i is not transitive. If j is any node of T_n not in E_i , and such a node exists since $m > 1$, then either $j \rightarrow E_i$ or $E_i \rightarrow j$. If $j \rightarrow E_i$ then $\Gamma(j)$ is non-empty and non-transitive. Thus the conditions are also necessary, in view of Lemma 1. This completes the proof of Lemma 2.

4. Proof of theorem 1. All highly-regular tournaments T_n have property \mathcal{L} , as was shown in [1], since they clearly satisfy the condition of Lemma 1. Thus

the tournaments described in the statement of Theorem 1 certainly have property \mathcal{L} , by Lemma 2.

Let T_n denote any tournament with property \mathcal{L} . We may assume that T_n is non-transitive and simple, in view of Lemma 2, and we may also suppose that $n \geq 3$. To complete the proof of Theorem 1 we must show that T_n is highly-regular.

The *score* of any node i is the number of nodes beaten by i . If s denotes the maximum of the scores of nodes of T_n , then $1 \leq \frac{1}{2}(n-1) \leq s$; furthermore, $s \leq n-2$ since if $s = n-1$ then T_n would be reducible and not simple. Let x denote any node with score s . We may suppose, for convenience, that x has label n ; that the nodes of $\Gamma(n)$ are labelled $1, 2, \dots, s$; and that the nodes of $\Gamma^{-1}(n)$ are labelled $s+1, \dots, n-1$. Since T_n has property \mathcal{L} it follows that $\Gamma(n)$ and $\Gamma^{-1}(n)$ are both transitive, by Lemma 1. So we may further assume that if $1 \leq i < j \leq s$, then $i \rightarrow j$ in $\Gamma(n)$; and that if $s+1 \leq u < v \leq n-1$, then $u \rightarrow v$ in $\Gamma^{-1}(n)$. This labelling defines a natural circular ordering of the nodes of T_n and, in what follows, when we refer to the successors or predecessors of a node, we mean the successors or predecessors with respect to this ordering. At this stage we may say that node n beats its first s successors (and loses to its $n-1-s$ predecessors) and that each node in $\Gamma(n)$ and $\Gamma^{-1}(n)$ beats its immediate successors in these subtournaments (and loses to its immediate predecessors in these subtournaments). We want to show that the same is true for each node i .

Suppose for some node x in $\Gamma(n)$ there exist two nodes u and v in $\Gamma^{-1}(n)$, where $u < v$, such that $u \rightarrow x$ and $x \rightarrow v$. Then $v \rightarrow n$, $n \rightarrow x$, and $x \rightarrow v$; that is, the nodes v , n , and x form a 3-cycle, and all three of these nodes lose to u . If this were the case, then $\Gamma(u)$ would be neither empty nor transitive and T_n would not have property \mathcal{L} . It follows, therefore, that if a node x in $\Gamma(n)$ beats any nodes of $\Gamma^{-1}(n)$, then those nodes form a subset of consecutive nodes of $\Gamma^{-1}(n)$ starting with node $s+1$.

Node 1 beats the $s-1$ nodes $2, \dots, s$ and, in addition to these nodes, it beats just one node of $\Gamma^{-1}(n)$. For, if node 1 beats more than one node of $\Gamma^{-1}(n)$, its score would exceed the maximum score s ; and, if node 1 lost to all the nodes of $\Gamma^{-1}(n)$, nodes 1 and n would be equivalent, contradicting the assumption that T_n is simple. Therefore, node 1 beats one node of $\Gamma^{-1}(n)$, namely, node $s+1$, and loses to the remaining nodes of $\Gamma^{-1}(n)$.

We next observe that node s must beat node $s+1$, for otherwise it would have score zero and T_n would be reducible and not simple. Finally, we assert that all remaining nodes of $\Gamma(n)$ must also beat node $s+1$. For, if there were some node x , where $1 < x < s$, such that $s+1 \rightarrow x$, the nodes x , s , and $s+1$ would form a 3-cycle in $\Gamma(1)$; then $\Gamma(1)$ would be neither empty nor transitive and T_n would not have property \mathcal{L} .

It follows from the preceding observations that node 1 beats its first s

successors (and loses to its $n-1-s$ predecessors) and that each node of $\Gamma(1)$ and $\Gamma^{-1}(1)$ beats its immediate successors in these subtournaments (and loses to its immediate predecessors in these subtournaments). By repeating this argument we find that the same is true for every node i of T_n . In particular, each node of T_n beats its s immediate successors. This implies that $s = \frac{1}{2}(n-1)$ and that n is odd. Hence, T_n is highly-regular by definition. This completes the proof of Theorem 1.

5. Enumerating tournaments with property \mathcal{L} . Let $f(n)$ denote the number of tournaments T_n with n labelled nodes that have property \mathcal{L} , and let $g(n)$ denote the corresponding number when the labels of the nodes are not taken into account.

THEOREM 2. *If $n = 1, 2, \dots$, then*

$$f(n) = (n-1)! 2^{n-1}$$

and

$$g(n) = \frac{1}{2n} \sum_k \phi(k) 2^{n/k}$$

where $\phi(k)$ denotes the Euler ϕ -function and the sum is over all odd divisors k of n .

Proof. Let $T_n = R_m(E_1, \dots, E_m)$ denote a tournament with property \mathcal{L} , where R_m is highly-regular. It follows from the definition of R_m that there exists a circular ordering of the subtournaments E_1, \dots, E_m such that $E_i \rightarrow E_j$ if and only if E_j is one of the first $\frac{1}{2}(m-1)$ successors of E_i with respect to this ordering (and the ordering with this property is unique). Furthermore, there is a unique linear ordering of the nodes in each transitive subtournament E_i such that each node u in E_i beats its successors with respect to the ordering in E_i . These orderings, that is, the circular ordering of the subtournaments E_i and the linear orderings of the nodes in the individual subtournaments E_i , induce a circular ordering of the n nodes of T_n such that each node u beats its immediate successors that belong to the same or one of the next $\frac{1}{2}(m-1)$ subtournaments E_i .

It follows, therefore, that there is a one-to-one correspondence between the labelled tournaments T_n with property \mathcal{L} and the circular arrangements of m 0's and the n numbers $1, 2, \dots, n$ such that no two 0's are next to each other. The numbers $1, 2, \dots, n$ correspond to the nodes of T_n and $i \rightarrow j$ in T_n if and only if there are at most $\frac{1}{2}(m-1)$ 0's between i and j in the circular arrangement. There are $(n-1)!$ circular arrangements of the numbers $1, 2, \dots, n$ and for each such arrangement there are $\binom{n}{m}$ ways to insert m 0's. Since m can be

any odd number not exceeding n , it follows that

$$f(n) = (n-1)! \left\{ \binom{n}{1} + \binom{n}{3} + \cdots \right\} = (n-1)! 2^{n-1},$$

as required.

Now suppose the labels of the nodes are not taken into account. It is not difficult to see that there is a one-to-one correspondence between the unlabelled tournaments $T_n = R_m(E_1, \dots, E_m)$ with property \mathcal{L} and the circular arrangements of n 1's and m 0's such that no two 0's are next to each other. (Two such arrangements are considered the same if they differ only by a rotation.) If no two 0's are next to each other, then each 0 is followed by a 1; thus, the circular arrangements of n 1's and m 0's with no two 0's next to each other are equinumerous with the circular arrangements of $(n-m)$ 1's and m 0's. It is well-known (see, e.g., [5; p.162]) that the number of such arrangements is

$$\frac{1}{n} \sum_k \phi(k) \binom{n/k}{m/k}$$

where the sum is over all divisors k of m and n . When we sum this expression over all odd numbers m not exceeding n we obtain the required formula for $g(n)$. (Since m is odd it follows that k is odd.)

Notice that when n is a power of 2, then

$$g(n) = \frac{1}{n} 2^{n-1};$$

and when n is an odd prime p , then

$$g(p) = \frac{1}{p} (2^{p-1} + p - 1).$$

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